Towards an Understanding of Scalable Query and Data Analysis for Social Media Data and Analysis on High-Level Dataflow Systems

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**ABSTRACT**

Nowadays there is great research potential in analyzing the vast amount of data collected from social media and social network applications. In order to explore the correlations among this data and social activities, modeling techniques such as data mining and machine learning solutions have been applied in combination with ad hoc query and complicated post-query data analysis. Use of high-level platforms such as Pig, Hive, and Spark SQL to support this type of sophisticated analysis has become the most popular solution. However, the question remains: which of the available software building blocks can serve users best according to their data needs? This question motivated us to research the execution flow and performance characteristics of these platforms, focusing on our special interests of social media data, to provide a detailed comparison for ad hoc queries and applications performed on top of them.

# INTRODUCTION

Social media data and its applications have gained the attention from commercial, academic and research communities. Many interesting research applications [1-6] that deal with daily activities, events, and knowledge in human society have been developed. The data collected by these metrics of social media is vast, far exceeding constraints found in the storage substate, databases, and runtimes traditionally used to store and access historical data. In practice, social media service providers such as Twitter, Facebook and Instagram have accommodated their users with in-house solutions. However, for those research scientists and application developers who subscribe to public social streams to build their research systems and prototypes, it is challenging to select appropriate software building blocks that can scalably store, serve, and customize data schema for such immense data.

Gao et al. [7-10] working with the *Truthy* [11] project illustrates the usefulness of Apache software stacks with *IndexedHBase* [12]. This in turn led to Truthy’s current deployment on a large-scale and large-storage private cluster, MOE. Though IndexedHBase and its Java API have met the fundamental requirements for accessing and processing data analysis, there are still unfulfilled areas where further research is viable, especially when integrating the existing analysis pipelines with Apache high-level language platforms such as Pig [13], Hive [14] and Spark SQL [15]. Other challenges include ad hoc queries and direct computation on top of storage and databases.

The scope of our research is outlined here in very fine-grained low-level perspective such as I/O(s) consumption benchmark, state-of-the-art system-level building block comparison [16], and possible performance optimization research in regards to different types of data processing. Its aim is to understand the requisite background knowledge, perform benchmarks, and review the existing research. Based on benchmark results we can identify the execution and performance characteristics for queries and applications run on these high-level platforms. Using these categories, we might further investigate the differences between social media data and general data analysis to identify the potential customizations for social data analysis on high-level platforms.

This paper focuses on understanding the requirements and boundaries of data systems that support various applications mixed with ad hoc queries and data analysis, especially social media data analysis. We benchmark query systems including Pig, Hive and Spark SQL. In particular, our previous work has investigated the possibility of using User Defined Functions (UDF) to support complicated iterative algorithms with fine-grained data aggregation and communication patterns [17]. We are working to construct and research the state-of-the-art end-to-end data pipeline for general scientific data and social media data.

The paper is structured as follows. Section 2 introduces our target data and data model, system-level requirements, and baselines in supporting social media data. Section 3 describes the features and characteristics of ad hoc queries, and our research target of using high-level languages with NoSQL databases for query and data analysis. Section 4 benchmarks applications discussed in Section 3. Lastly, we draw our conclusion and summarize the research directions in Section 5 and Section 6.

# TRUTHY - SOCIAL MEDIA OBSERVATORY

*Truthy* is a public social media observatory developed as a research project at Indiana University. It analyzes and visualizes information diffusion on Twitter. Truthy monitors and collects Twitter data in real-time directly through the Twitter public steaming API [18]. Much of our work has been done recently with the help of this observatory, and researchers have yielded important inferences of human society by analyzing the social activities in cyberspace.

One example of end-to-end social media data analysis [2] involved utilizing the IndexedHBase queries [12] on top of further data mining techniques, such as eigenvector modularity [19] and label propagation [20]. The analysis was carried out on two datasets about political discussion collected during the six weeks leading up to the 2010 U.S. congressional midterm elections and 2012 U.S. presidential elections. The results shown in [2, 12] prove that the retweet networks exhibited a highly segregated partisan structure; users of those tweets are mainly split into two homogenous communities corresponding to the political left and right leanings. Figure 1 shows the execution flow for getting the graph of political polarization. In 2012, the average amount of collected tweets each month was about 1 billion tweets. Such a huge dataset proved problematic in terms of storage, as was providing a fast processing layer to handle such large amounts of data.

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| Figure 1. Sophisticated pipeline for visualizing Political Polarization |

This data observatory has been storing Twitter streaming data since July 2010; the current data size as of Aug 2015 is approximately 162TB. It includes raw compacted JSON files on HDFS, tweet fields and inverted indices stored as HBase tables. IndexedHBase [12] has been used to create inverted indices for raw tweets in JSON format. Various fields such as keywords, hashtags, geographical locations, user IDs, and retweet IDs have been stored as searchable rowkeys while the related tweet IDs are stored as associated (multi-column) values. Scientists and developers of Truthy perform ad hoc queries and post-query data analysis on these HBase tables, where tweet tables (JSON fields) and index tables are semi-structured with different amounts of columns.

## System Challenges for Truthy

We have compared different NoSQL solutions to support indexing and fast queries on large-scale social media data. As a result, IndexedHBase was selected as the framework to store, serve, and perform data computation for data scientists [7-10] using YARN Hadoop and HBase as building blocks.

Utilizing an infrastructure supported by IndexedHBase, our work delves into the system challenges for ad hoc queries and post-query data analysis performed on large-scale social media data. Based on our studies, the three categories of challenges are data-related, system-related, and programming and computation-related. Data-related challenges store and serve incremental data on a scale of at least TB level, sustain or create indices with customized formats, and support flexible data schema for structured and semi-structured data with less disk consumption. System-related challenges offer multi-tenancy to query clients and application developers, as well as allowing commodity hardware failures with fast/auto recovery. Finally, programming and computation-related challenges support ad hoc query interface such as Pig, Hive, and Spark SQL, in addition to supporting customized programming in imperative programming languages such as Java and Python. They offer different levels of parallelism and sophisticated data mining and machine learning applications.

# AD HOC QUERY WITH NOSQL DATABASE

A key characteristic of social media data analysis is the ad hoc queries that select the interested subset of data from a very large set of time spatial data stored in databases. Generally, each row/field of tweet data is stored with an associated timestamp and their related column values. An example query could be “Find all the related tweets with given hashtag #computing in the time range between June 15th 2015 and July 10th 2015”. This type of query can be rewritten as traditional Select-Project-Join (SPJ) ad hoc queries. These project and join the two datasets, the records within that specific time, and other sets of records within the target fields, such as hashtags. The size of projection data, amount of generated temporary tables, and the type of join operations depends on the target fields of each query within a single table. For instance, the execution flow of the example query given above firstly scans the entire raw data table and filters the required data by referring to the given predicates of time duration and a hashtag. Then it generates two temporary tables and performs a single shared-key join. Due to the extra overhead of generating two tables separately, in addition to performing a join aggregation and scanning entire rows of each target record, Gao et al. [7-9] has shown that the overall performance does not meet our expectations. By comparison, the HBase solution scans the index and raw tables once and immediately filters the data with the support of built-in “create timestamp” for each stored row/column in a table. Even adopting NoSQL databases as backend storage, there are limited choices of database solutions that can efficiently store large datasets with fast (inverted) index access to the time spatial data. IndexedHBase was developed as the backend inverted index layer, where the data and indices are stored on top of HBase to support these complicated social media data queries.

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| **Type** | **Query** | **Exe. Steps** |
| Read-One-Write-One | get-tweets-with-meme, get-tweets-with-text, get-tweets-with-userid, get-retweets, get-tweets-with-time, get-tweets-with-phrase | 2 |
| Read-One-Transform-One | get-retweet-edges, get-mention-edges | 2 |
| Read-One-Transform-Many | meme-post-count, text-post-count, userid-post-count, user-post-count, user-post-count-by-text, meme-cooccur-count | 2 |
| Single-Scan | meme-timestamp-count, text-timestamp-count, userid-timestamp-count | 1 |

Table 1. Classification of support social queries

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| Figure 2. Dataflow for Ad hoc queries of social media data |

Most of our support queries are HBase I/O intensive, which mainly perform random data access by specified row keys, e.g. tweet IDs to tweet table and keywords to text index tables. Each query must first retrieve the related tweet IDs from index tables by a given time range and queried keys. It then obtains the required columns from the tweet table and may perform a UDF to yield a stage-ready result output on HDFS for further data analysis as shown in Figure 2. This differs from SQL database procedure. IndexedHBase must build the indices as separate tables on HBase, and it considers extra overheads when loading data into HBase. Based on the execution flow and different type of data transformation of these queries, we have identified four categories of supported queries as shown in Table 1:

1. *Read-One-Write-One*: Obtain one related tweet ID from Index Table by the given queried key (e.g. hashtag), dump the whole tweet as result, e.g. get-tweets-with-meme.
2. *Read-One-Transform-One*: Obtain one related tweet ID from Index Table by the given queried key (e.g. hashtag), generate single output entry (e.g. user pair) from the obtained tweet, e.g. get-retweet-edges.
3. *Read-One-Transform-Many*: Obtain one related tweet ID from Index Table by the given queried key (e.g. hashtag), generate multiple output entries as ArrayList, e.g. meme-cooccur-count.
4. *Single-Scan*: read the statistic information directly from HBase table.

## Query Execution with High-level Languages

IndexedHBase includes Java MapReduce implementations driven by a wrapper bash shell. Despite this, it is not easy to add new queries or UDF without understanding the background of Hadoop MapReduce. Specifically, all the supported social media data queries are very straightforward ad hoc queries executed with common database operations such as *FILTER*, *GROUP BY*, *JOIN*, and FOR EACH with built-in or UDF functions. This motivated us to investigate the integration with high-level abstractions such as *Pig* [13], *Hive* [14], and *Spark SQL* [15] for day-to-day query and data analysis.

Most of these systems are considered as *Dataflow* system or *Dataflow programming* model, which is a paradigm that models a program as a directed graph of data [21]. In both cases, data flows among a series of components such as operators and functions which serve as a “*black-box*” unit (the detailed implementations are already defined) to transform the incoming data from its original format into another. Data in the execution flow is clearly defined as either being input or output to every atomic component, independently handled on each and inherently run in parallel.

*Pig* [13] is a dataflow system built on top of Hadoop MapReduce, which aims to serve as a high level abstraction interfacing with SQL database and MapReduce computation systems. Pig itself is a declarative DAG-flow system, but it uses *Pig-Latin* [22], a procedural language. This makes it flexible and allows users to choose different implementations of the same relational operator (e.g. *JOIN* and *GROUPBY*) in execution. Other than the built-in operators, a developer can apply their own sophisticated algorithm to the dataflow in Pig via its UDFs. *Hive* [14] is another high-level platform, but it differs from Pig by supporting data warehouse ad hoc queries and simple MapReduce applications for structured data stored on *HDFS* [23]. It provides a SQL-like language, *HiveQL*, to execute on top of Hadoop. Most of the implementation concepts of Hive derive from SQL RDBMS. *Spark SQL* [15] is another open source project inspired by Hive. Instead of being tightly coupled with the Hadoop MapReduce engine, it uses Spark as its low-level runtime, with DataFrame schema RDD as its major in-memory data structure embedded with named column (table-like) schema. The extensible query optimizer Catalyst is written in Scala, a different model from Hive and its predecessor *Shark* [24].

We will discuss the overhead of using these high-level platforms for the target ad hoc queries in Section 4. In addition, although we have not yet linked the ad hoc query with the post-query analysis, we recognize the need for chaining this intermediate data to next-generation compute resources and fulfilling the dataflow of the entire analysis pipeline. Our previous work [17] has demonstrated the importance of in-memory computation and resource reuse for sophisticated machine learning applications with iterations. We not only incorporate the Hadoop plugin Harp, but allow general ETL queries to continue the sophisticated application immediately afterward [15]. This would save significant job restart overhead and enable fast resource allocation and reusability. Furthermore, it enables intuitive development writing prototypes of end-to-end pipelines in a single environment. Spark SQL has proposed a similar idea that uses the same platform and data abstractions for both queries and analysis, yielding meaningful results for sophisticated algorithms. Meanwhile, *Apache Tez* [25] shows the importance of resource reusability for complex DAG tasks on top of high-level platforms run on YARN Hadoop.

# PERFORMANCE RESULT

Our experiments run on MOE, a large-storage, large-memory and high-performance private cluster at Indiana University devoted to the Truthy project [26, 11]. It consists of 3 login nodes and 10 compute nodes, where each login node is set up with two Intel(R) Xeon(R) CPU E5-2620 v2 CPUs, 64 GB memory, and each compute node has five Intel(R) Xeon(R) CPU E5-2660 v2 CPUs, 128 GB memory, 48TB HDD and 120GB SSD. All nodes are interconnected with a 10Gb Ethernet. We perform our tests on top of a Hadoop 2.5.1 cluster with different high-level platforms such as Pig 0.14.0, Hive 1.0.0, and Spark SQL 1.5.0. Meanwhile, IndexedHBase 0.2.0 is the Java MapReduce baseline.

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| Figure 3. Performance breakdown for get-tweets |

We have implemented a total of 17 ad hoc queries [12] written in all the platforms mentioned above. Other than the initial stage of searching related tweet IDs from index tables, we compare these three query implementations on different platforms and examine their runtime behaviors as shown in Figure 3, Figure 4, and Figure 5. Each submitted query runs with a total of 587858 tweet IDs obtained from meme index tables by being given the most common hashtag, *“Follow Friday”* *#ff* and are equally assigned to 9 workers; the default parallelism (the amount of reducers) is set to 4.

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| Figure 4. Performance breakdown for get-retweet-edges |
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| Figure 5. Performance breakdown for meme-cooccur-count |

Since all of these are HBase I/O intensive queries, the main brunt of overhead is the data retrieval time communicated with the HBase tweet table which stores the original tweet fields. Other than get-tweets query, which dumps the entire tweet to HDFS, every implemented UDF only scans a subset of columns and yields a specified format such as edge pair (user ID and retweet user ID) and a list of mentioned hashtags in the related tweet. These transformed data are collected and accumulated by using the standard data aggregation operations, e.g. *GROUP BY* and *reduceByKey*. Compared with traditional row-based databases, we save significant I/O overhead with the help of the columnar scanning provided by HBase. Note that the computation time of Spark SQL takes longer as Spark performs “map-only” worker execution; it includes the UDF transformation time (from DataFrame RDD to Java RDD) and cross-worker data aggregation/communication time, along with the output to HDFS time.

In addition, since all these queries are compiled and run as YARN or Hadoop jobs, we also evaluate the local write bytes (except Spark SQL which does not have a reduce stage) and investigate the data aggregation overhead. As shown in Figure 6, for queries with reduce stages (*get-retweet-edges* and *meme-cooccur-count*), Pig and Hive implementations have more intermediate data and match the trend of overall execution time. This is due to these high level abstractions using tuple-based computation and emitting each processed tuple to the output buffer. There IndexedHBase is pure Java MapReduce implementation with which the output of the mapper is optimized, combining the emitted values that shared the same output key. We also observe this behavior from the intermediate record sizes as shown in Table 2.

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| Figure 6. Intermediate Local write in bytes |
| **Query** | **Pig** | **Hive** | **IndexedHBase** |
| get-tweets | 587858 | 587858 | 587858 |
| get-retweet-edges | 179486 | 179463 | 167740 |
| meme-cooccur-count | 90216 | 90125 | 63524 |

Table 2. Mapper output (combined if any) record sizes

# CONCLUSION

This paper compares social media data query performance on a large-scale data observatory. By addressing the challenges in various levels of this observatory, we proposed the use of inverted indices generated by IndexedHBase with different high-level abstractions to perform query analysis and post-query data analysis. Our argument here and in future work is that the simplest solution offers the greatest potential. By this we mean programming interface, computation extension, and data linkage should be constructed within a single platform. Doing so could achieve better resource utilization by reducing the resource allocation overhead, fast data access with in-memory caches for frequently used data within a pipeline, and even better query execution flow by referring to the real-time and statistical data metrics of the processing data.

As may have been observed, our research does not investigate the *query optimization* of databases [27-30] with optimization strategies such as predicates move-around [29], which have been implemented in many database [31-33, 14] and dataflow [13] systems, especially for Select-Project-Join (SPJ) ad hoc queries. However, as mentioned above, the social media data queries can prove challenging for traditional SPJ database systems. Our implementation therefore bypasses the SPJ complexity by using inverted indices with associated timestamps within the same cell of data.

# FUTURE WORK

We have integrated Harp with Pig [17] to show the advantages of using customized data aggregation and in-memory computation for iterative applications. We also wrote a high-level comparison survey [16] to qualify the basic features and fundamental differences among Pig, Hive and Spark SQL. Based on these efforts, we plan to extend our research direction with a quantitative understanding of the state-of-the-art Apache high-level language platforms for end-to-end solutions that link multiple compute components into a single development and platform. We will also revisit the behaviors of running the computations on these high-level platforms versus domain-specific languages such as R and Matlab.

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