# AIRCRAFT DESIGN PROBLEM IMPLEMENTATION UNDER THE COMMON OBJECT REQUEST BROKER ARCHITECTURE

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# <u>Abstract</u>

The paper describes a component-based computational environment for implementing aircraft design problems. The environment is conducive to taking advanta ge of the paral lelisms inher ent in the problem and distribute the individual disciplines on machines most appropriate to their needs while insulating the developer and the user from the complexity of the underly ing commu nications const ructs. Commo n Objec t Reque st Broker Archi tecture and Java programming language are used to encap sulate discipline codes as "objects". An interface file identifies all the information needed by a user of the object. Legacy codes are "wrap ped" using Java's native interface methodology and each such code is called from a modul e which implements the services of the object. A server program ties the implementation to the interface. Data and file management are accomplished using Java's datab ase connectivity to access a commercial relational database manag ement syste m. Java's Beans Devel opment Kit is used to implement the disciplines and sub-tasks as reusable compo nents that provide a graph ical interface for user input as well as facilitate interactive and visual object connectivity and problem execution progress monit oring. This approach has been used to implement a simpl e aircr aft desig n optimizat ion probl em, the analy sis part of a large scale high speed civil trans port desig n optim ization problem, and a stand alone aerodynamic optim izer.

#### **Introduction**

Conventional aircr aft desig n requires diver se engineering disciplines to execute independent of each other, in sequence, and often times iteratively and interactively. This process is further complicated by the fact that the focus, emphasis, and approach of each discipline can be quite distinct, and multiple invocations of the discipline programs are often required to arrive at a feasible design. The end result is a set of thumbprint, carpet, or, correlation plots from which a "best" design may be chosen. The whole process is vastly time-consuming and does not, in general, include all engineering disciplines early in the design process [1,2].

Preli minary desig n tradi tionally deals with disciplina ry sizin g and shapi ng, with relia nce on previ ous desig ns of a vehic le type. The vehic le must susta in several criti cal flight (load ing) condi tions throu ghout the opera tional envel ope durin g which the loads are typically redis tributed due to aeroe lastic effects. Analytic al and test verif ication of desig ns may be performed throu ghout the desig n process [1]. Other conditions such as flutt er, diver gence, control revers al, and gust loadi ng must also be addre ssed. Howev er, these conditions may not be included in more detail until static stren gth desig n is completed [3,4]. Information from each discipline is analyzed and the desig n is modif ied in a seque ntial, itera tive proce ss.

The tradi tional appro ach resul ts in a desig n procedure that is large ly infle xible and computationally taxing. New techn iques in multi disciplina ry desig n optimization are aimed at improving desig n efficiency, design cycle time reduction, and introducing decision critical information early in the desig n process [5]. An earlier effort within the Frame work for Inter disciplina ry Design and Optimization (FIDO) project [6] used Parallel Virtu al Machine (PVM) to handle communications among vario us discipline codes executing in a "host/slave" mode. This frame work was sensi tive to the host operating system and changing the analytical connectivity or switching discipline codes required major programming inter vention.

The goal of the current frame work is to provide a programming environment for automating the distribution of a complex computing task over a networked, heter ogeneous syste m of computers. These computers inclu de engin eering works tations, vecto r may super computers, and paral lel processing computers. The present paper describes a computational environment for multi disciplina ry analy sis and optim ization of a High Speed Civil Trans port (HSCT) aircr aft, capable of concurrent analyses in a distributed computing environment and significantly reducing the design cycle time while intro ducing more detailed analysis early in the desig n process.

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The present approach incor porates advanced aircraft design techniques within a new design frame work utili zing the Commo n Objec t Reque st Broke r Architecture (CORBA) and the Java programming language. The prima ry benefit of this system is the flexi bility to demon strate advan ced techn ology concurrent multi discipli nary desig n integ ration techn iques and the capab ility to introduce detailed analyses early in the design proce ss. Such a multi disciplina ry analy sis and optim izatio n (MDO) syste m capable of concurrent analyses using several disciplines such as aerodynamics, structures, missi on/perform ance and optimization is under development for the High Performance Computation and Communication and Computational Aeros ciences (HPCC P/CAS) program. This syste m uses CORBA in an Object Frame work [7] for the integrated design of a HSCT aircr aft confi guration acros s a netwo rked syste m of heter ogeneous computers using the client-server parad igm. A central relational datab ase is used for informa tion inter change and file management. The goal of the design is to minimize the vehicle gross takeoff weight (GTOW) for given flight conditions and mission require ments.

The paper will describe two implementations in the aircr aft desig n domai n based on the objec t-oriented appro ach using CORBA and Java. The first is a simple repre sentative aircr aft desig n problem based on fast limit ed-fidelit y disci pline codes inclu ding an equivalent plate struc tural analy sis, linea r aerod ynamic analysis, table lookup for propulsion and a simple range equation for performance fuel weight estimation. The second implementation is of a stand alone aerodynamic optimizer based on high-fidelity discipline codes includi ng a nonli near aerod ynamics code, param etrized geome try codes and an optim ization code. The components of the techn ologies used, CORBA and Java, will be described and their use in the implementation will be detai led. Compo nents devel oped for the second impleme ntation were reused without change in the Analysis part of a large scale desig n of a HSCT aircr aft.

### **The Framework**

A multi disciplina ry analy sis and optim ization syste m has been devel oped that is capab le of concu rrent analy ses using sever al disci plines such as aerod ynamics, struc tures, performance, propulsion, and optim ization. This syste m, based on CORBA and the Java programming langu age and its Appli cations Programming Inter face's (API)'s, uses the client-server paradigm in an Object Frame work for imple menting problems such as the integ rated design of a high speed civil trans port aircr aft over a networked syste m of heter ogeneous computers. The Beans Devel opment Kit (BDK) is used to provi de a Java Beans -based graph ical interface for user input, inter active and visual object conne ctivit y, and for monit oring the progress of problem execution. Java's Datab ase Conne ctivity (JDBC) is used by the clien t and serve r objec ts to communicate with a central relational datab ase. Java's Remot e Method Invocation (RMI) is used on platforms where a costeffective commercial implementation of CORBA is not available. In the three applications described in this paper, optimization plays a common role. In the case of the representative high speed civil trans port desig n problem, the objec tive is to optimize the airplane weight for given cruis e conditions, range, and paylo ad requirements, subject to aerod ynamic, structural, and performance constraints. The desig n variables include both structural thick ness and geometric parameters defining the airplane shape. The frame work provides the capability to switch between low-, medium-, and high-fidelity codes with ease.

# <u>Common Object Request Broker Architecture</u> (CORBA):

The Commo n Objec t Reque st Broke r Archi tecture is a speci fication adopt ed by a conso rtium of indus try repre sentatives known as the Objec t Manag ement Group (OMG) to define a frame work for devel oping distr ibuted appli cations. CORBA allow s clien t objec ts to invok e server objec ts acros s the network witho ut havin g to deal with the under lying compl exities of object imple mentation and invoc ation. In this model, an objec t is an encap sulated entit y with a uniqu e ident ity whose services can be accessed only through a well defined *inter face*. The imple mentation of the objec t (lang uage, opera ting syste m, other syste m specific aspects) as well as the locat ion of the invok ed objec t are trans parent to the reque sting clien t. The detai ls of the archi tecture are discu ssed below.

CORBA can be thought of as a "soft ware bus" [Fig. 1] connecting various objects, both application and service, on a network of computers. Objects on the bus can be used by any other objects on the bus, with the Object Request Broker (ORB) mediating the transfer of messages between them. In this configuration, there is peer-to-peer communication where servers can be clients for the services of other objects on the bus. CORBA also defines a wide range of services and facilities [8] to extend the core capabilities of the ORB.



Figure 1. CORBA as a 'software bus'.



Figure 2. The Object Request Broker, ORB.

## **Objec t Reque st Broke r (ORB) :**

A softw are imple mentation of the CORBA specificat ion is called the Object Reque st Broker, or, ORB [Fig. 2]. The ORB media tes the trans fer of messa ges from a program to an object on a remote networked host. The ORB deliv ers reque sts to objec ts and returns any responses. The key feature of the ORB is the transparen cy of how it facil itates the clien t/object commu nication. The client is not required to know where the target object resides, how and in what programming langu age it was implemented, or the operating system on the host computer. When a client makes a request, the client is not concerned wheth er that object is currently active and ready to accept requests. The ORB trans parently activates the object, if required, before delivering the request. The client does not need to know what under lying communication mechanism the ORB uses to media te the messa ge passi ng betwe en the clien t and the serve r. All these enable the user to generate "thin clients" i.e., all the number crunching is done on the server-side on computers most appropriate for the task. The ORB frees the appli cation developer to focus more on the application domain issues and less about the low-level distributed system programming issue s.

An ORB is one component of the OMG's Object Management Architecture (OMA). The other s include the application objects, CORBA services, and CORBA facilities. Services include a) Naming Service - which allows the clients to find objects based on names, b) Trading Service - which allows clients to find objects based on their properties. CORBA facilities define a

set of high-level services that applications frequently require when manipulating distributed objects.

Diffe rent comme rcial imple mentations of the ORB must all be able to talk to each other using a stand ard netwo rk proto col calle d the Inter net Inter -ORB Protocol (IIOP).

Within an object frame work, each object communicat es with other s on a peer- to-peer basis. Each object is a client of other servi ces and a server f for the servi ces it provi des. Very often, a client for one request is a server for anoth er. This architecture facil itates network programming by allow ing the creation of distributed appli cations as sets of coope rating reusa ble objects that inter act as though they were implemented in a single programming langu age on one computer.

### Interface Definition Language (IDL):

Altho ugh CORBA objec ts are implemented using stand ard progr amming langu ages, each CORBA object has a clear ly defined interface, specified in the CORBA Inter face Defin ition Languag e. Before a client can make a request to an object, the client must know the types of operations supported by the object. An object's interface specifies the operations and types that the object supports and thus defines the requests that can be made to that object. These object interfaces, writt en in IDL, are simil ar to class es in C++ and to interfa ces in Java. IDL is a decla rative langu age, not a programming langu age. It forces interfaces to be separate from object implementations. To use a discipline code, the user needs to know only what interfaces are imple mented by the code.

The IDL interface is compiled by an IDL Compiler. IDL compilers are available for several programming langu ages. For Java, the IDL compiler produces several Java constructs which correspond to the IDL definition. The mapped constructors may be divided into those that allow a client to access an object through the object interface, and another set of constructs which allow the object to be implemented in a server.

### Java and its APIs:

Java is a general-purpose concurrent class-based objectoriented programming language that fits naturally into the CORBA object orientation architecture. Java is specifically designed to have as few implementation dependencies as possible and allows application developers to write a program once and then be able to run it everywhere on the Internet. Java is robust, architecture neutral, portable and with the "just-in-time" compilers, approaches speeds comparable to those of languages such as C or FORTRAN. Java's multithreading capability provides the ability to execute multiple activities in parallel. Java's many APIs such as RMI, JNI, JDBC and Java Beans (disc ussed below), make Java the programming langu age of choic e for distributed applications, parti cularly those dependent on legacy codes, relational datab ases and requiring a graph ical user interface.

**Remote Method Invocation (RMI):** Remote Method Invocation (RMI) is Java's alternative to CORBA's remote invocation capabilities and is designed to simplify the communication between two objects residing on different hosts. RMI is useful only for communication between Java objects, and it does not currently use a stand ard transmission proto col such as IIOP. This API is a handy alternative when a suitable commercial ORB is not available (or too expensive) for a specific platform. Of course, the platform should have the Java Virtu al Machine (JVM) on it.

Java Native Interface (JNI): While a pure Java solution is nice in principle, realistically, for an application such as airplane design, there are several situations where it becomes necessary to use codes written in another language. Java's Native Interface methodology permits calling such legacy codes from a Java program. To make calling native methods possible, Java comes with hooks for working with system libraries and a few tools to relieve some of the associated tedium. However, the usage of native methods precludes portability. In the present framework, the portability issue would not be a problem since these methods would be used on the server side of the application.

Java Database Connectivity (JDBC): The JDBC API is a set of specifications that define how a program written in Java can communicate and interact with a database. JDBC defines how the communication is to be carried out and how the application and database interact with each other. More specifically, the JDBC API defines how an application opens a connection, communicates with a database, executes SQL statements, and retrieves query results. JDBC provides a vehicle for the exchange of SQL between Java applications and databases. JDBC classes are available for several popular commercial databases which makes it possible to switch databases on an application without being required to make any significant changes to the Java code.

**JavaBeans:** A JavaBean is a reusable software component that can be manipulated visually in a builder tool. The builder tools may include web page builders, visual application builders, GUI layout builders, or even server application builders. The JavaBeans API provides an environment in which a programmer can "wrap" an object as a component that can be used by other developers. JavaBeans provide the capability and a set of standards for a design or builder tool to be able to query the

component package and access the object's properties through a process known as introspection.

JavaB eans have three distinct elements: Properties, Event s, and Metho ds. Properties are the inter nal varia bles associated with a component. Event s are a way for components and applications to communicate with each other. Common event s include mouse movements and click s, keys being pressed, objects receiving or losing focus, etc. In addition, the user can add custo m bean event s to handle special requirements. Metho ds are functions that the component can perform by invoc ation from the outsi de world.

JavaB eans can expos e selec ted properties for a user to set at design time or get at run time. Properties can be as simple as a file name or as sophi sticated as a color editor or arrays of data. Simple properties are displayed in the JavaB ean's Property Sheet, while more sophi sticated properties require custom built property editors accessible from the bean's Property Sheet. Customizers can also be written to permit users to edit multiple properties at the same time and make this graph ical user interaction more user friendly.

### **Object** Creation:

In the curre nt imple mentation, Iona Techn ologies' imple mentation of the CORBA stand ard for the Java programming langu age, Orbix Web, was chose n as the ORB. For each disci pline, to be wrapped as an object, an interface file is written in IDL (Java for RMI implementation) identifying the services offered by that object, the required input parameters, the outputs, and the types of errors the object can "throw". Figure 3 shows a listing of an IDL interface for the implementation of the airpl ane drag estimation discipline code as "aeroDrag" object. The inter face for aeroDrag has only one method, indicating that the object provides only one service - getTo talDrag that takes as input two float ing point value s: cruis e angle of attack and the press ure drag coefficient, and returns a float ing point value, the total drag coefficient - cd\_to tal..

The object and its services, as identified in the interface file, are then implemented in Java in an *implementation* file. The services are often obtained by the execution of a discipline code. Discipline codes are in general legacy codes, written either in FORTRAN or C, and are accessed through an intermediate function which is created following the guidelines set in JNI for calls to

module Aerodynamics
{
// Exception checks if error code is returned
exception gotNegativeFlag
{
long errorNo;
};
interface aeroDrag {
//Operations
float getTotalDrag( in float alfa_cruise,
in float cd_pres )
raises (gotNegativeFlag);
};
};



native funct ions. Within the implementation file, a centr al relat ional datab ase is queri ed for neede d data and file infor mation. Any required file management is done based on this file infor mation. A third item of softw are associated with the creat ion of a disci pline object is the "Serv er" class code. This component ties the implementation class to its IDL interface. Serve rs provi de objects for use by clients and other serve rs.

In a client code (the fourth of four files associated with an object), the service *getTotalDrag* from object *aeroDrag* is obtained by first instantiating the object as

# aeroDrag my\_aeroDrag = aeroDragHelper.bind(":aeroDragSrv","cmb");

where *my\_aeroDrag* is an instance of *aeroDrag* and the implementation is linked to the interface by the object reference *aeroDragSrv*. The service is being requested of a server on the computer "cmb". The actual request for service is done by:

# float cd\_total = my\_aeroDrag.getTotalDrag (alfaCruise, cd\_pres);

To create a JavaBean, the client code is "wrapped" following the guidelines set by Java's Beans Development Kit. This wrapper is associated with a BeansInfo file which identifies the discipline and other properties exposed to the user through a Property Sheet and associated custom property editors. Customizers, if any, would be implemented in separate software components.

A distr ibuted appli cation can then be assembled in one of two ways. In the first approach, a 'mast er' client program can be writt en to implement the desig n

analysis algor ithm by making service calls to the distributed application objects. In the second approach, the custom discipline JavaBeans can be imported into either BDK's Beanbox, or into one of several commercially available application builders such as Java Studio, Visual Cafe', Jbuilder, etc. Once imported, these beans become available as icons or menu options. A user would select appropriate beans from the menu and place them in the work area. A "connection wizard" is usually used to connect the beans together to form the required analytical network. The connections are made by tying events to methods. One type of connection handles sequencing the executions of the various discipline beans. Another type of event called the Property Change Event manages the communication of changes in exposed properties between discipline beans dependent on these properties. In this environment, graphical user interface component beans can be integrated into the application including buttons, labels, text windows and chart beans to monitor execution progress. The application execution can be started by a simple clicking on a button.

In either case, when the master program is executed or when the button is pressed, the client calls are transferred to the ORB which then passes the function calls through the server code to the target object. Components of the ORB are implemented by the OrbixWeb daemons running on the server hosts or by the RMIregistries in the case of Remote Method Invocation.

# <u>Simple Airplane Design Optimization</u> <u>Problem</u>

This section describes the implementation of a simplified airplane design optimization benchmark problem in the CORBA/Java based Object Frame work. This problem is considered an excellent multidisciplinary optimization test case since the interplay of multiple disciplines is attempted while carry ing along only a small number of design variables, constraints, and a single objective function. Figure 4 shows the example HSCT model, without engine masses and control surfaces, that was studied within this frame work. Figure 5 presents the discipline segments and the information flow necessary for the design and analysis of an optimal HSCT configuration.

The princ ipal disci plines for the desig n problem are: aerod ynamics, struc tures, propulsion, and performance. The desig n objec tive is to minim ize the aircr aft gross take-off weight for a given cruise conditions, range, and paylo ad requirements. The weight is minimized subject to aerod ynamic, struc tural, and performance constraints such as the limit ing values of lift and drag, maxim um stres ses at critical points on the wing inboard and outboard panels, and range. Desig n vari



Figure 4. HSCT model problem and design point data

ables include wing sweep, wing root chord, distance to wing sweep angle break, and the inboard and outboard skin thickness.



Figure 5. Flow-chart of Design Optimization Problem

After initialization, the design optimization proceeds in three phases, a) analysis, b) gradient computation, and c) optimization. The analysis phase begins with a calculation of drag polars using the mediumfidelity code Wingdes [9]. Lift and drag values for a range of angles of attack are used in generating parametric representations of aerodynamic responses. All subsequent aerodynamic analyses for that design cycle will utilize these drag polars to compute lift and drag. The next step in the analysis phase is the iteration for the airplane weight convergence. The weight iteration loop begins with a static trim analysis where force balance is computed for two different load factors: a) load factor = 1.0 for drag calculation used in performance analysis, and b) load factor = 2.5 for loads calculation in structural design. The propulsion segment computes the current fuel flow rate. The performance segment uses this flow rate to produce an estimate for fuel weight.

A structural analysis is done once during the first iteration to determine the structural weight. A loads transfer program converts the aerodynamic pressure distribution over the airplane to vertical forces on the structure at a trimmed angle-of-attack and a load factor of 2.5. The structural analysis program used here is the Equivalent Laminated Plates Solution, ELAPS [10]. The total weight is then computed as the sum of the fixed weights, structural weight, and the fuel weight. The process is repeated until the total weight converges within a predefined tolerance.

In the next phase, all the system response derivatives required by the optimizer are computed. The gradi ents of aerod ynamic and struc tural const raints are computed using finit e-differences. Gradi ents of the fuel weight with respect to design variables are obtained using a close d-form expression.



Figure 6. JavaBeans Connectivity.

In the third phase, Conmin [11] with linear approximations is used as the optimization program. Conminuses the method of useab le-feasible directions to minimize the objective function, the airplane gross weight, subject to the aforementioned design constraints and computes an updated set of values for the design variables.

The problem disciplines were wrapped as objects in the CORBA/Java environment. The client side codes were wrapped as JavaB eans and imported into BDK's Beanbox. Figure 6 shows the analytical connectivity of the discipline JavaB eans in the Beanbox. The Design Optimization Problem is presented in Fig. 5.



Figure 7. Variation of Structural Design Variables.



Figure 8. Variation of Airplane Weights.

A typic al HSCT confi guration flying at Mach 2.4 at an altit ude of 63,000 feet was analyzed using the analy tical conne ctivity shown in Fig. 5. Figure 7 shows the varia tion of the struc tural design varia bles with cycle number while Fig. 8 shows the varia tion of the airpl ane weight components with cycle number. The results show excel lent agreement with results obtaine d from an imple mentation of the current problem in an earli er PVM based framework[12].

## Stand-alone Aerodynamic Optmizer

A second appli cation implemented in the CORBA /Java distr ibuted computing frame work is a stand -alone aerod ynamic optim izer. The appli cation uses the paral lel computation of aerod ynamic deriv atives via autom atic differentiation of the Euler /Navier-Stoke s solve r CFL3D [13] coupled with an optimizer and surfa ce/volume grid deformation tools to perform an optim ization to reduce the drag of a HSCT airpl ane configuration.

Centr al to any gradi ent-based problem is the evalu ation of solut ion deriv atives with respect to the chose n desig n varia bles. Differe ntiation of the CFD source code used to obtain the solut ion gives exact deriv atives of the discrete equat ions, without the step size problems of finite differences[14]. A parameterized surface definition is used that relates the shape to geometric design varia bles. The method is a free-form deformation approach very simil ar to morphing technique s used in computer anima tion.

Figure 9 shows the analytical connectivity for the application. The design cycle begins with grid generation. Grid generation includes the following steps: a) use the updated design variables along with the parame-



Figure 9. Stand-alone aerodynamic Optimizer.

teriz ed grid to form an updat ed surfa ce grid and its sensi tivity deriv atives with respect to the desig n variables, b) use the updat ed surfa ce grid to gener ate a volum e grid and corre sponding sensi tivity deriv atives, and c) convert the volum e grid and sensi tivity data to binar y forma t. The former may be divid ed from a singl e block to multi ple block s so that CFL3D may be execute d in a parallel mode

Grid gener ation is follo wed by the analy sis phase in which an analy sis is conducted using CFL3D. This can be follo wed by a gradi ent computation phase. During this phase a flow solver analy sis is follo wed by sensi tivity analy sis. Under assum ed HSCT flight conditions the flow solution converges rapid ly and hence the analy sis phase is skipp ed in favor of a combined analysis and gradi ent calculation phase.

In the next phase, the optim izer computes new values of the desig n varia bles. The updat ed desig n varia bles are the input to the next cycle in the optim izatio n process. The overall analysis is run for a preset number of desig n cycles. The unconstrained optim izatio n problem uses 27 shape desig n varia bles to maxim ize the value of the lift coefficient.

The analysis codes which include the geometry analysis codes, CFL3D, and the optimizer were wrapped as objects in the CORBA/Java framework. The clientside was wrapped using JavaBeans technology to provide a graphical user interface through which the user is



Figure 10. Variation of  $C_D$  with cycle number.

able to modif y selec ted input conditions including options to start the analysis from scratch or continue from a selected cycle in a previous design analysis. Norma lized drag results for an optimization run of 10 cycles is presented in Fig. 10. Both scaled and unscaled design variables are shown for the current implementation. These results compare identically with those obtained by the conventional process.

## A High Fidelity Aircraft Design Problem

The experience gained from the earlier implementations and a number of the objects developed therein are being used in the CORBA/Java implementation of a large scale aircraft design optimization problem titled HSCT4.0. The objective is to demonstrate the application of high performance computing techniques to the problem of multidisciplinary optimization of a subsonic transport configuration using high fidelity analysis simulations early in the design process. The HSCT4.0 problem considers 27 shape design variables and 244 structural design variables (the 0/45/90/core layers in 61 optimization zones). Figure 4 shows the example model and Fig. 11 shows a high level flow chart of the multidisciplinary analysis portion of the design optimization problem.

During the Analysis phase, the values of the objective function and the constraints are evaluated. The process begins by deriving an updated geometry and corresponding linear, nonlinear aerodynamic meshes and a structural finite element mesh by applying the updated shape design variables to a baseline parameterized geometry grid.

Next, in the Weights process, structural weight from a finite element analysis, fuel weight from a performance analysis, and empirical weights representing all other components such as passengers, seats, actuators, etc. are all evaluated and assembled to provide the



Figur e 11. Flowc hart of analy sis in MDO probl em.

"as built" weight in terms of nodal weights and a total weight and center of gravity location for cruise and take off conditions.

A Nonlinear Corrections module periodically computes "corrections" for use with subsequent linear aerodynamics computations. The process uses updated linear and nonlinear aerodynamics grids, aircraft total weight for each load condition considered, along with multiple runs of the linear aerodynamics code covering a range of bracketing angles-of-attack and tail angles of incidence. For each load condition, the angle-of-attack is determined that generates the same configuration normal force as that of the nonlinear aerodynamic calculation. The nonlinear corrections are calculated as the difference between the Z components of the linear panel loads, interpolated at the angle of attack that matches the normal force, with the nonlinear panel loads that have been transferred to the linear aero grid.

The Rigid Trim calculation for the cruise condition uses the derived linear aerodynamic grids along with the nonlinear corrections to compute the angle-ofattack and tail angle of incidence, that produces the target lift coefficient (as per the weight estimation), with no net pitching moment. The resulting surface pressures and the induced drag coefficient are then determined by the final trim angle and tail angle of incidence.

Once the Rigid Trim process has been completed, the Polars, Performance, and Ground Scrape sequence of processes may proceed in parallel with the Displacements, Loads Convergence, and the Stress & Buckling process sequence.

Aircraft drag polars are calculated for the current design over a range of Mach numbers, altitudes, and angles-of-attack so as to provide the required input to the performance module. At each mission condition, lift and drag due to lift are computed from linear aerodynamics calculations that span the Mach number range. Empirical corrections are applied to drag due to lift. In addition, wave drag and friction drag are computed to obtain the total drag.

The table of lift coefficients and coefficients of drag components at the mission points are input to the performance module, the outputs from which include range, take off field length, landing field length, lift off speed, approach speed and time to climb to cruise.

The ground scrape process provides a basis for constraints such that the aircraft tail will not scrape the ground on take off or landing. The ground scrape constraints are formulated as limits on the maximum value of the take off and landing gross weights for avoiding tail scrape. Additionally, ground clearances are computed for selected airframe and engine locations at the take off and landing angles of attack and a given pitch angle.

The trimmed cruise loads are input to the Displacements module. Within the Displacements module, the pressure loads on the aerodynamic grid are transferred as consistent loads at the nodes of the structural mesh. The cruise displacements obtained from a structural analysis are then input to the Loads Convergence module.

Converged loads for selected flight conditions are obtained through an aeroelastic analysis. At each load condition, the aircraft is trimmed for a consistent set of loads, angle of attack, and tail incidence angle. The trimmed loads are transferred to the structural nodes. Displacements computed from a structural analysis are then used to deform the aerodynamic grid. The aircraft is then trimmed at this new configuration and the trim loads transferred to the aircraft structure, followed by another structural analysis. This process is repeated till the trimmed loads (and the corresponding displacements) converge to a pre-set tolerance. These converged loads for the selected load conditions are input to the Stress and Buckling module. In this module, the Hoffman Stress Failure Index (SFI) is computed for every face sheet for each of 2260 elements for 7 flight conditions. A buckling load factor (BLF) is computed for each of the elements in terms of the inplane stress resultants.

All components of the Analysis part of the HSCT4.0 problem have been implemented as objects in the current CORBA/Java environment. Some objects developed for the earlier examples, such as the Geome try and Nonli near Corrections objects from the Stand alone Aerod ynamic Optimizer, were used without modification.

## **Conclusion**

A compo nent-based computational environment for imple menting aircr aft design problems has been described. The environment is conducive to taking advantage of the paral lelisms inherent in the problem and distribute the individual disciplines on machines most appropriate to their needs while insulating the devel oper and the user from the complexity of the underlying communications constructs. CORBA and the Java programming langu age are used to encap sulate discipline codes as "objects".

Under the CORBA a multi discipline analysis and optimization system has been developed that is capable of concurrent analysis. The primary disciplines considered in this paper are aerodynamics and structures coupled with a formal optimization technique. All objects are implemented using Java, with each CORBA object having a clear ly defined interface. Java is a gener al-purpose concurrent class -based object orien ted progr amming langu age. Java is robus t, architectu re neutr al, porta ble and with just- in-time compi lers, approaches speeds comparable to those of language s such as C and FORTR AN. While a pure Java solut ion is nice in princ iple, reali stically, for an application such as aircraft design it becomes necessary to use codes writt en in other langu ages. This is permitted throu gh Java's Native Interface Techniques, which permits codes written in other languages to be called from a Java program. The Java JDBC API allows Java programs to communicate and interact with a central database. JavaB eans have been used to implement the objects as reusable software components that permit user interaction and visual object connectivit y. JavaB eans allow the user to selec t properties at design time or access at run time.

Three desig n probl ems were presented. The first example is the "Simple Airpl ane Desig n Probl em" which is multi discip linary in nature. The contributing disci plines are aerod ynamics, propulsion, performance and structures coupled with a formal optim ization techn ique. Results show excellent agree ment with earli er bench mark results. The second example is the "Stan d-alone Aerod ynamic Optim izer". Results of the optim ization compared ident ically with those obtained by the conventional process. Good design convergence was achieved within ten design cycles. Final ly the implementation of the analysis portion of a large scale high speed civil trans port design optim ization problem is described.

A component-based multi disciplina ry analysis and optim ization design frame work that will allow global coupling of several uncoupled engineering disciplines has been demonstrated using the CORBA. Contributing disciplines function concurrently as objects in the frame work. Design trends compared well with results obtained from conventional technique s.

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